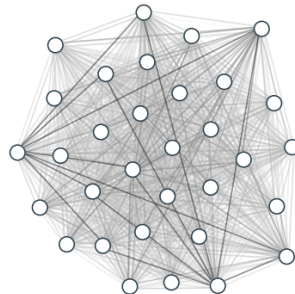


# A Topological Centrality Measure for Directed Networks

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# Motivation & Background

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Networks model complex systems as (directed) graphs

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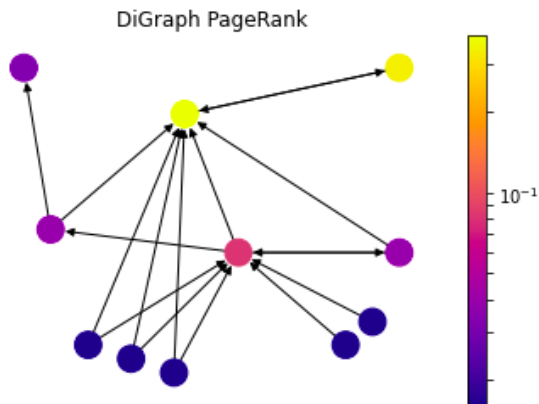
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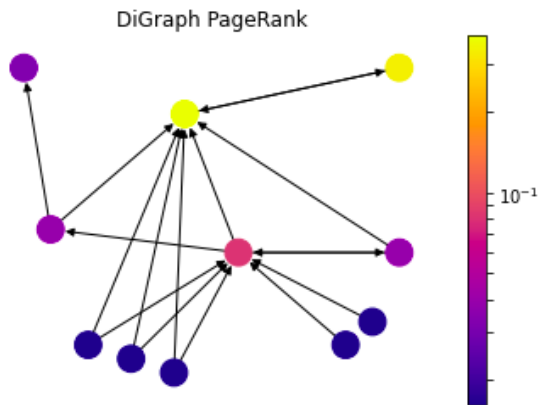


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## Node Centrality

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## Goal

Define a centrality measure that captures non-local propagating effects and directedness.

# Networks



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## Definition (F.Iannelli, 2017)

Let  $G = (X, w_X)$  be a network, define  $\gamma(G)$  to be  $(X, m_X)$  where  $m_X : X \times X \rightarrow \mathbb{R}$  is given by:

$$m(x, y) = \begin{cases} 1 - \log \frac{w(x, y)}{\sum_{z \neq y} w(x, z)} \geq 1 & \text{if } y \neq x \\ 0 & \text{if } y = x \end{cases}$$

Two nodes that interact a lot ( $w(x, y) \gg 0$ ) will be closer ( $m(x, y) \sim 1$ ).

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Given a network  $G$  and  $x$  a node in  $G$ , define  $f(G, x) = (X \setminus \{x\}, w_X|_{X \setminus \{x\}})$ , i.e. the sub-network induced by deleting  $x$  and all edges incident to  $x$  in  $G$ .

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## Idea

Given  $x$  a node in  $G$ , we compare the difference in the “[dis]connectivity” of  $\gamma(G)$  and  $\gamma(f(G, x))$ .

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We use the "size" of the homology of a "shape" built from  $G$  as a proxy for disconnectivity.



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Recall a *simplicial complex* is a set of tetrahedrons of any dimension "glued together in a nice way".

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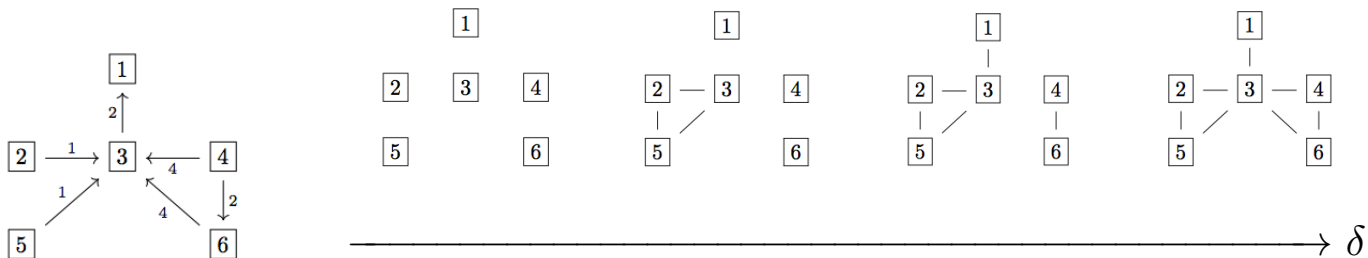
## Definition (F.Memoli and S.Chowdhury, 2016)

Given a network  $G = (X, w_X)$  and  $\delta \in \mathbb{R}$ , the **Dowker Complex**  $\mathcal{D}_{\delta, G}$  is the simplicial complex given by:

$$\mathcal{D}_{\delta, G} := \{\sigma \subseteq X : \exists p \in X \text{ s.t. } w(x, p) \leq \delta \forall x \in \sigma\}.$$

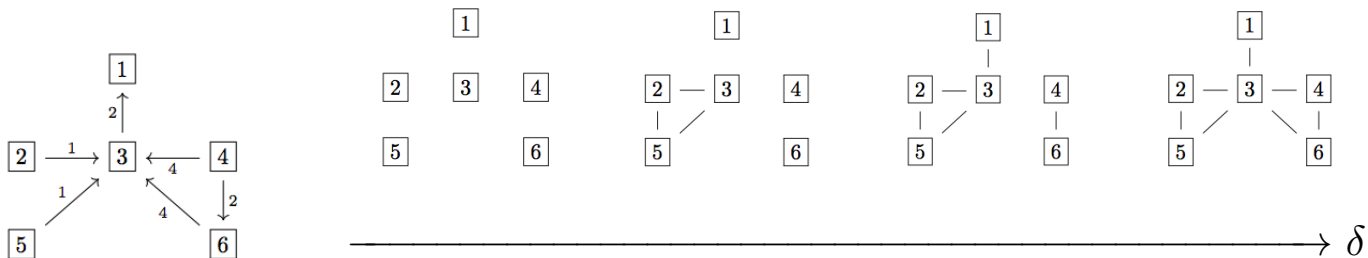
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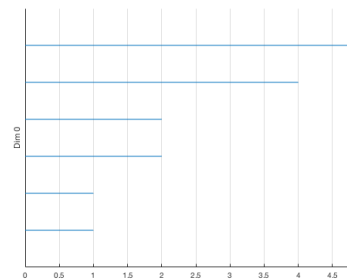


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- 1 As  $\delta \nearrow$ , number of path components  $\searrow$ .
- 2 This data is recorded on a **persistence diagram**.
- 3 We denote  $\mathbf{P}_0(G)$  as the set of 0-dimensional barcodes for the Dowker complex  $\mathcal{D}_{\cdot, G}$ .



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## Definition

Let  $G$  be a network. The **quasi-centrality**  $C(x)$  for node  $x \in X$  is:

$$C(x) = \sum_{c \in \mathbf{P}_0(f(\gamma(G), x))} \text{length}(c) - \sum_{c \in \mathbf{P}_0(\gamma(G))} \text{length}(c) + d$$

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## Theorem

For a network  $G = (X, w_X)$ ,  $C(x)$  is nonnegative for all  $x \in X$ .

# Applications

## Goals

- Demonstrate that  $C$  is a valid measure of centrality
- Use quasi-centrality to assess the influence of a node in a real-world network.



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- Interdependency between far-flung communities
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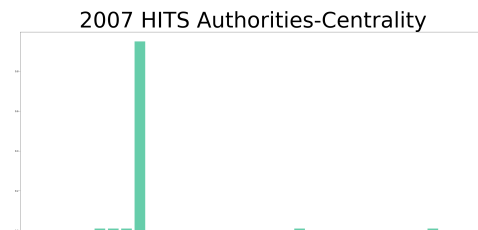
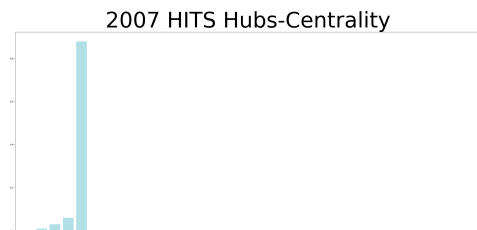
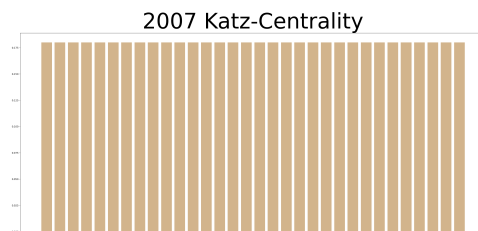
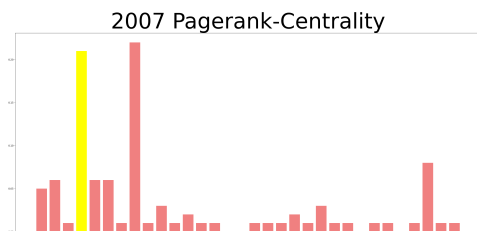
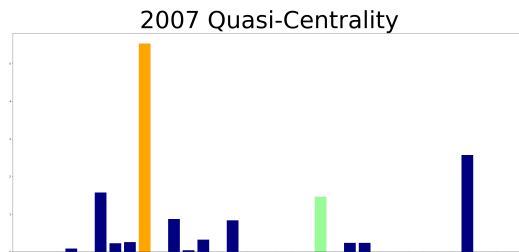
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## Data

- OECD Inter-Country Input-Output (ICIO) Tables
- Machinery production network in Asia
- Industries: machinery equipment, computer and electronics, electrical machinery, auto machinery

# Results



# Future directions

- Compute the quasi-centrality measure for other asymmetric networks
  - ▶ biological networks
  - ▶ airflight networks
- Relate higher dimensional homological features in directed networks to real-world phenomena
  - ▶ trade flows
  - ▶ embargo
- Define other measures in network analysis using TDA
  - ▶ connectivity
  - ▶ robustness
  - ▶ efficiency

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