

Coin Games and 5-way Scales

Joshua Lee

Mentor: Dr. Tanya Khovanova

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Three Way Scale



- N total identical-looking coins
- Balance scale
- 1 fake, lighter than real
- Goal: determine fake coin using the least amount of weighings as possible.

Information Theoretical Bound

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- $w = \lceil \log_3 N \rceil$

Five Way Scale

- Two more possible outcomes
- $d = \#$ of fake coins on left pan - $\#$ of fake coins on right pan

MUCH LESS	LESS	EQUAL	MORE	MUCH MORE
$d \geq 2$	$d = 1$	$d = 0$	$d = -1$	$d \leq -2$

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Lemma

After one weighing, let a_1, a_2, a_3, a_4, a_5 be the number of remaining possibilities of the fake coins for the outcomes MUCH LESS, LESS, EQUAL, MORE, MUCH MORE respectively. Then,

$$\max(a_1, a_2, a_3, a_4, a_5) = a_3$$

regardless of how many coins were on each pan.

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- $\frac{3}{5} \cdot 5^w \geq \binom{N}{2}$

- Linear Strategy — Compare 2 coins with 2 coins: $w = \left\lfloor \frac{N}{2} \right\rfloor$

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- Better Strategy — Divide into 3 equal piles
 - ① MUCH LESS: Problem Reduced to $\frac{N}{3}$
 - ② LESS: $w = \log_2 N$
 - ③ EQUAL: One more weighing to reduce problem to $\frac{N}{3}$
- $w = 2 \lceil \log_3 \frac{N}{3} \rceil$ weighings

- Better Strategy

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- Improved Bounds

Conjecture

If w denotes the maximum number of weighings in any strategy that guarantees finding the fake coins, and N is the total number of coins, then there exists a constant k such that

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for all N .

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- General n-way Scale

- Nim — Basic Game



- Goal: Take the last stone

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- Goal: Take the last stone
- Winning (P) positions and losing (N) positions

Minimum Excluded

Definition

The **minimum excluded value** (often shorted as **mex**) of a subset of some well-ordered set is the smallest value not included in the set.

For our use, we will assume that we are using set of non-negative integers.

Example

- $\text{mex}(0, 1, 3) = 2$
- $\text{mex}(1, 2, 3) = 0$
- $\text{mex}(0, 2, 4, 6\dots) = 1$

Sprague-Grundy Theorem

Every impartial game is equivalent to a nim-heap of a certain size.

- Game equivalent to the number of stones in Nim
- mex of the set of reachable Grundy Numbers

The following research was begun by the following people:

- Kyle Burke
- Tanya Khovanova
- Richard J. Nowakowski
- Amelia Rowland
- Craig Tennenhouse

- Aequitas — Latin concept of equity
- Game regarding the classic coin problem
- Must reveal information every turn
- Observer cannot know the fake coin
- Player loses if there is no legal move

Grundy Numbers for Aequitas

- One Final position - 2 remaining possible coins
- P-position - final position
- N-positions - every other position

Game Values

N	Grundy Number
$4k + 3$	$2k$
$4k + 4$	$2k + 1$
$4k + 5$	$2k + 1$
$4k + 6$	$2k + 1$

Modified Aequitas; Game 2

- Fake coin either heavier or lighter
- Observer cannot know fake coin
- One Final position — only P-position

Game Values

N	Grundy Number
$2k$	$2k - 2$
$2k + 1$	1

Modified Aequitas; Game 3

- Observer cannot know relative weight of fake coin
- Two final positions - only P-positions
- Equal Grundy numbers as Game 2

Game Values

N	Grundy Number
$2k$	$2k - 2$
$2k + 1$	1

Future Research

- Limit to number of coins on each scale
- Other games

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